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# A Lower Bound on Independence in Terms of Degrees

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## Abstract

We prove a new lower bound on the independence number of a simple connected graph in terms of its degrees.

**Keywords:** Independence; stability; connected graph

**AMS subject classification:** 05C69

## 1 Introduction

We consider *finite, simple, and undirected graphs*  $G$  with *vertex set*  $V$ . For a graph  $G$ , we denote its *order* by  $n$  and its *size* by  $m$ , respectively. The *degree* of  $u$  in  $G$  is denoted by  $d(u)$  and  $\Delta$  is the *maximum degree* of  $G$ . A set of vertices  $I \subseteq V$  in a graph  $G$  is *independent*, if no two vertices in  $I$  are adjacent. The *independence number*  $\alpha$  of  $G$  is the maximum cardinality of an independent set of  $G$ . The independence number is one of the most fundamental and well-studied graph parameters [6]. In view of its computational hardness [5] various bounds on the independence number have been proposed. The following classical bound holds for every graph  $G$  and is due to Caro and Wei [1, 7]

$$\alpha \geq \sum_{u \in V} \frac{1}{d(u) + 1}. \quad (1)$$

Since the only graphs for which (1) is best-possible are the disjoint unions of cliques, additional structural assumptions excluding these graphs allow improvements of (1). A natural candidate for such assumptions is connectivity.

For connected graphs, Harant and Rautenbach proved [2] (cf. also [3] and [4])

**Theorem 1** *If  $G$  is a connected graph, then there exist a positive integer  $k \in \mathbb{N}$  and a function  $\phi : V \rightarrow \mathbb{N}_0$  with non-negative integer values such that  $\phi(u) \leq d(u)$  for all  $u \in V$ ,*

$$\alpha \geq k \geq \sum_{u \in V} \frac{1}{d(u) + 1 - \phi(u)}, \quad (2)$$

and

$$\sum_{u \in V} \phi(u) \geq 2(k - 1). \quad (3)$$

Note that Theorem 1 is best-possible for the connected graphs which arise by adding bridges to disjoint unions of cliques, i.e. it is best-possible for the intuitively most natural candidate of a connected graph with small independence number. In [3], a weaker version of Theorem 1 is proved. This result is obtained from Theorem 1 by replacing the inequality (3) by  $\sum_{u \in V} \phi(u) \geq k - 1$ .

For an integer  $l$  with  $0 \leq l \leq 2m$  let  $f(l) = \min \sum_{u \in V} \frac{1}{d(u)+1-\phi(u)}$ , where the minimum is taken over all integers  $0 \leq \phi(u) \leq d(u)$  with  $\sum_{u \in V} \phi(u) = l$ .

Obviously,  $f$  is strictly increasing. With this function  $f$ , it follows the existence of positive integers  $k_1$  and  $k_2$  such that  $\alpha \geq k_1 + 1 \geq f(k_1)$  (put  $k_1 = k - 1$  and use the result in [3]) and  $\alpha \geq \frac{k_2}{2} + 1 \geq f(k_2)$  (with  $k_2 = 2(k - 1)$  and Theorem 1). After extending  $f$  to real arguments, in [4], it is proved that the function  $l + 1 - f(l)$  is continuous and strictly increasing and that  $k_1$  is at least the unique zero  $k_0$  of this function. Finally,  $\alpha \geq k_0 + 1$  is the main result in [4].

Here we will show that the continuous function  $\frac{l}{2} + 1 - f(l)$  is also strictly increasing. If we assume that  $\sum_{u \in V} \frac{1}{d(u)+1} \geq 2$  for the graph in question then  $f(2) > f(0) = \sum_{u \in V} \frac{1}{d(u)+1} \geq 2$ . It will be proved that there is a unique solution  $l_0$  of the equation  $\frac{l}{2} + 1 = f(l)$  and, because  $\frac{l}{2} + 1 - f(l)$  is strictly increasing and  $\frac{2}{2} + 1 - f(2) < 0$ , it follows  $l_0 > 2$ . Consequently,  $\frac{l_0}{2} + 1 = f(l_0) > f(\frac{l_0}{2} + 1)$ , since  $f$  is strictly increasing, hence,  $\frac{l_0}{2} > k_0$ .

The inequality  $\alpha \geq \frac{l_0}{2} + 1$  is the content of the following Theorem 2.

In case  $\sum_{u \in V} \frac{1}{d(u)+1} \geq 2$ , Remark 2 gives a lower bound on the improvement

$$(\frac{l_0}{2} + 1) - (k_0 + 1) = f(l_0) - f(k_0).$$

**Theorem 2** *Let  $G$  be a finite, simple, connected, and non-complete graph on  $n \geq 3$  vertices of size  $m \geq n$ . Moreover, let  $\alpha \leq \frac{n}{2}$  be the independence number,  $\Delta$  be the maximum degree of  $G$ ,  $n_j$  be the number of vertices of degree  $j$  in  $G$ , and*

$$x(j) = \frac{j(j+1)}{j(j+1)-2} \left( \left( \frac{2}{j+1} - (\Delta - j) \right) n_\Delta + \dots + \left( \frac{2}{j+1} - 1 \right) n_{j+1} + \frac{2n_j}{j+1} + \dots + \frac{2n_1}{2} - 2 \right)$$

for  $j \in \{2, \dots, \Delta\}$ .

Then

(i) *there is a unique  $j_0 \in \{2, \dots, \Delta\}$  such that  $0 \leq x(j_0) < n_\Delta + \dots + n_{j_0}$  and*

(ii)

$$\begin{aligned} \alpha &\geq \left( \sum_{u \in V} \frac{1}{d(u)+1} \right) + \frac{n_\Delta}{\Delta(\Delta+1)} + \frac{n_\Delta + n_{\Delta-1}}{(\Delta-1)\Delta} + \dots + \frac{n_\Delta + \dots + n_{j_0+1}}{(j_0+2)(j_0+1)} + \frac{x(j_0)}{(j_0+1)j_0} \\ &= \frac{x(j_0) + n_{j_0+1} + 2n_{j_0+2} + \dots + (\Delta - j_0)n_\Delta}{2} + 1. \end{aligned}$$

## 2 Proof of Theorem 2

In the sequel let  $k$  be the lower bound on  $\alpha$  of Theorem 1.

By Theorem 1, it follows

**Lemma 1**  $k \geq f(2(k-1))$ .

For a finite family  $F$  of integers let  $\max(F)$  be a maximum member of  $F$ . Note that a member of a family may occur more than once. If for instance  $F = \{1, 2, 2\}$  then  $(F \setminus \{\max(F)\}) \cup \{\max(F) - 1\} = \{1, 1, 2\}$ . The following Lemma 2, Lemma 3, and Lemma 4 are proved in [3] and [4].

**Lemma 2** *Given an integer  $l$  with  $0 \leq l \leq 2m$ , the following algorithm calculates  $f(l)$ :*

*Input: The family  $F = \{d(u) \mid u \in V\}$ .*

*$j := 0$ ,*

*while  $j < l$  do begin  $F := (F \setminus \{\max(F)\}) \cup \{\max(F) - 1\}$ ;  $j := j + 1$  end*

*Output:  $f(l) = \sum_{m \in F} \frac{1}{m+1}$ .*

Lemma 3 is a consequence of Lemma 2.

**Lemma 3** *Given an integer  $0 \leq l \leq 2m$ ,*

*(i) there are unique integers  $j$  and  $x$  with  $j \in \{1, \dots, \Delta\}$  and  $x \in \{0, \dots, n_\Delta + \dots + n_j - 1\}$  such that*

$$l = n_\Delta + (n_\Delta + n_{\Delta-1}) + \dots + (n_\Delta + n_{\Delta-1} + \dots + n_{j+1}) + x$$

$$= x + n_{j+1} + 2n_{j+2} + \dots + (\Delta - j)n_\Delta$$

*and*

$$(ii) f(l) = (n_\Delta + \dots + n_j - x) \frac{1}{j+1} + \frac{x}{j} + \frac{n_{j-1}}{j} + \dots + \frac{n_1}{2}$$

$$= (n_\Delta + \dots + n_j) \frac{1}{j+1} + \frac{x}{j(j+1)} + \frac{n_{j-1}}{j} + \dots + \frac{n_1}{2}.$$

By Lemma 3, it follows

**Lemma 4** *If  $l = x + n_{j+1} + 2n_{j+2} + \dots + (\Delta - j)n_\Delta$  with  $j \in \{1, \dots, \Delta\}$  and*

*$x \in \{0, \dots, n_\Delta + \dots + n_j - 1\}$  then  $f(l+1) - f(l) = \frac{1}{j(j+1)}$ .*

Using Lemma 3, the calculation of  $f(l)$  is possible without taking a minimum and without using the algorithm above. We will now define the function  $f$  for real  $l \in [1, m]$ . For given  $j \in \{1, \dots, \Delta\}$  and a real number  $x$  with  $0 \leq x < n_\Delta + \dots + n_j$  let the real numbers  $l$  and  $f(l)$  (implicitly) be defined as  $l = x + n_{j+1} + 2n_{j+2} + \dots + (\Delta - j)n_\Delta$  and  $f(l) = (n_\Delta + \dots + n_j) \frac{1}{j+1} + \frac{x}{j(j+1)} + \frac{n_{j-1}}{j} + \dots + \frac{n_1}{2}$ .

We will prove Lemma 5.

**Lemma 5** *The function  $g$  with  $g(l) = \frac{l}{2} + 1 - f(l)$  is continuous and strictly increasing on  $[1, n]$ .*

**Proof.** Consider  $l \in [1, n]$ . Then there are  $j \in \{1, \dots, \Delta\}$  and  $x$  with  $0 \leq x < n_\Delta + \dots + n_j$  such that  $l = x + n_{j+1} + 2n_{j+2} + \dots + (\Delta - j)n_\Delta$ .

If  $j = 1$  then  $n > l \geq n_2 + 2n_3 + \dots + (\Delta - 1)n_\Delta = 2m - n$ , a contradiction to  $n \leq m$ . Hence,  $j \geq 2$ , and  $l$  belongs to the interval

$$I(j) = [n_{j+1} + 2n_{j+2} + \dots + (\Delta - j)n_\Delta, n_j + 2n_{j+1} + \dots + (\Delta - j + 1)n_\Delta).$$

By Lemma 3,  $g(l + \epsilon) - g(l) = \epsilon(\frac{1}{2} - \frac{1}{j(j+1)})$  and, consequently,  $g(l)$  is continuous and, because  $j \geq 2$ , strictly increasing on  $I(j)$ .

Note that  $I(j) \cap I(j') = \emptyset$  if  $j \neq j'$  and that  $I(2) \cup \dots \cup I(\Delta) = [1, 2m - n] \supseteq [1, n]$ .

It is easy to see that  $g$  is also continuous in  $l = n_{j+1} + 2n_{j+2} + \dots + (\Delta - j)n_\Delta$  for  $j \in \{3, \dots, \Delta - 1\}$ . □

Since the classical bound due to Caro and Wei is tight only for complete graphs, it follows

$g(0) = 1 - \sum_{u \in V} \frac{1}{d(u)+1} < 0$ , and, by Lemma 1,  $g(2(k-1)) \geq 0$ . Using Lemma 5, there is a unique zero  $l_0 = x(j_0) + n_{j_0+1} + 2n_{j_0+2} + \dots + (\Delta - j_0)n_\Delta$  of  $g$  with  $1 < l_0 \leq 2(k-1) \leq 2(\alpha-1) < n$  and  $0 \leq x(j_0) < n_\Delta + \dots + n_{j_0}$ . It follows Lemma 6.

**Lemma 6**  $\alpha \geq k \geq \frac{l_0}{2} + 1$ , where  $l_0 \in (0, n]$  is the unique solution of  $\frac{l}{2} + 1 = f(l)$ .

Considering the equation  $\frac{l_0}{2} + 1 = f(l_0)$ , i.e.

$$2 + x(j_0) + n_{j_0+1} + 2n_{j_0+2} + \dots + (\Delta - j_0)n_\Delta = 2((n_\Delta + \dots + n_{j_0})\frac{1}{j_0+1} + \frac{x(j_0)}{j_0(j_0+1)} + \frac{n_{j_0-1}}{j_0} + \dots + \frac{n_1}{2})$$

it follows

$$x_0 = \frac{j_0(j_0+1)}{j_0(j_0+1)-2} \left( \left( \frac{2}{j_0+1} - (\Delta - j_0) \right) n_\Delta + \dots + \left( \frac{2}{j_0+1} - 1 \right) n_{j_0+1} + \frac{2n_{j_0}}{j_0+1} + \dots + \frac{2n_1}{2} - 2 \right).$$

We obtain Lemma 7.

**Lemma 7** *If  $j \in \{2, \dots, \Delta\}$  and  $l = x + n_{j+1} + 2n_{j+2} + \dots + (\Delta - j)n_\Delta$  with  $0 \leq x < n_\Delta + \dots + n_j$ , then  $\frac{l}{2} + 1 = f(l)$  if and only if*

$$x = \frac{j(j+1)}{j(j+1)-2} \left( \left( \frac{2}{j+1} - (\Delta - j) \right) n_\Delta + \dots + \left( \frac{2}{j+1} - 1 \right) n_{j+1} + \frac{2n_j}{j+1} + \dots + \frac{2n_1}{2} - 2 \right).$$

Now we complete the proof of Theorem 2. By Lemma 4 and Lemma 6,

$$\alpha \geq k \geq f(l_0) = f(0) + (f(1) - f(0)) + \dots + (f(\lfloor l_0 \rfloor) - f(\lfloor l_0 \rfloor - 1)) + (f(l_0) - f(\lfloor l_0 \rfloor))$$

$$= \left( \sum_{u \in V} \frac{1}{d(u)+1} \right) + \frac{n_\Delta}{\Delta(\Delta+1)} + \frac{n_\Delta + n_{\Delta-1}}{(\Delta-1)\Delta} + \dots + \frac{n_\Delta + \dots + n_{j_0+1}}{(j_0+2)(j_0+1)} + \frac{x(j_0)}{(j_0+1)j_0} \text{ because}$$

$$l_0 = x(j_0) + n_{j_0+1} + 2n_{j_0+2} + \dots + (\Delta - j_0)n_\Delta$$

$$= n_\Delta + (n_\Delta + n_{\Delta-1}) + \dots + (n_\Delta + n_{\Delta-1} + \dots + n_{j_0+1}) + x(j_0).$$

With  $f(l_0) = \frac{l_0}{2} + 1 = \frac{x(j_0) + n_{j_0+1} + 2n_{j_0+2} + \dots + (\Delta - j_0)n_\Delta}{2} + 1$  Theorem 2 is proved.  $\square$

### 3 Remarks

The following Remark 1 is proved in the introduction.

**Remark 1** *If  $\sum_{u \in V} \frac{1}{d(u)+1} \geq 2$  then  $\frac{l_0}{2} > k_0$ .*

Remark 2 compares the lower bound  $\frac{l_0}{2} + 1$  on  $\alpha$  in Theorem 2 to the lower bound  $k_0 + 1$  on  $\alpha$  in the main result in [4].

**Remark 2** *If  $\sum_{u \in V} \frac{1}{d(u)+1} \geq 2$  and*

$$k_0 = n_\Delta + (n_\Delta + n_{\Delta-1}) + \dots + (n_\Delta + n_{\Delta-1} + \dots + n_{j+1}) + x$$

*with  $0 \leq x < n_\Delta + \dots + n_j$  then  $\frac{l_0}{2} - k_0 \geq \frac{k_0}{j(j+1)}$ .*

**Proof.** Remark 1 implies  $\frac{l_0}{2} - k_0 = f(l_0) - f(k_0) > f(2k_0) - f(k_0)$ .

According to Lemma 2, the family  $F$  contains the member 1, the member 2, ..., and the member  $\Delta$  exactly  $n_1$  times,  $n_2$  times, ...,  $n_\Delta$  times, respectively.

Therefore, let the output  $f(l)$  of the algorithm in Lemma 2 be denoted by  $f_{n_1, \dots, n_\Delta}(l)$ .

With this notation, for example  $f_{n_1, \dots, n_\Delta}(1) = f_{n_1, \dots, n_{\Delta-1}+1, n_\Delta-1}(0)$ .

Using Lemma 3 (ii), it follows  $f_{n_1, \dots, n_\Delta}(k_0) = f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(0)$  and

$$f_{n_1, \dots, n_\Delta}(2k_0) = f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(k_0).$$

Consequently,

$$\begin{aligned} f_{n_1, \dots, n_\Delta}(2k_0) - f_{n_1, \dots, n_\Delta}(k_0) &= f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(k_0) - f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(0) \\ &= (f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(k_0) - f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(k_0 - 1)) \\ &\quad + (f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(k_0 - 1) - f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(k_0 - 2)) + \dots \\ &\quad + (f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(1) - f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(0)). \end{aligned}$$

Note that the expressions  $f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(s) - f_{n_1, \dots, n_{j-2}, n_{j-1}+x, n_\Delta + \dots + n_{j-x}}(s-1)$  equal fractions of type  $\frac{1}{a(a+1)}$  (see Lemma 3 and Lemma 4) with  $a \leq j$  for  $s = 1, \dots, k_0$ .

Thus,  $f_{n_1, \dots, n_\Delta}(2k_0) - f_{n_1, \dots, n_\Delta}(k_0) \geq \frac{k_0}{j(j+1)}$ .  $\square$

For integers  $r \geq 2$  and  $s \geq 2$ , consider the graph  $G_{r,s}$  obtained from  $s$  copies of the clique  $K_r$  on  $r$  vertices and adding  $s - 1$  mutually independent edges between these cliques such that  $G_{r,s}$  is connected. It follows  $\Delta = r$ ,  $n_j = 0$  for  $j < r - 1$ ,  $n_{r-1} = sr - 2(s - 1)$ ,  $n_r = 2(s - 1)$ , and  $\alpha = s$  for  $G_{r,s}$ . Using Theorem 2, we obtain

$$x(r - 1) = \frac{(r-1)r}{(r-1)r-2} \left( \left( \frac{2}{r} - 1 \right) n_r + \frac{2n_{r-1}}{r} - 2 \right) = \frac{(r-1)r}{(r-1)r-2} \left( \left( \frac{2}{r} - 1 \right) 2(s - 1) + \frac{2sr-4(s-1)}{r} - 2 \right) = 0 .$$

Hence,  $j_0 = r - 1$ ,

$$\frac{x(j_0) + n_{j_0+1} + 2n_{j_0+2} + \dots + (\Delta - j_0)n_\Delta}{2} + 1 = \frac{n_r}{2} + 1 = s = \alpha, \text{ and Remark 3 follows.}$$

**Remark 3** *There are infinitely many graphs  $G$  such that the lower bound on  $\alpha$  of Theorem 2 is tight.*

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